

Internetwork Protocols

Background to IP
IP, and related protocols

Internetworking Terms (1)

- Communications Network
 - Facility that provides data transfer service
- An internet
 - Collection of communications networks interconnected by bridges and/or routers
- The Internet - note upper case I
 - The global collection of thousands of individual machines and networks
- Intranet
 - Corporate internet operating within the organization
 - Uses Internet (TCP/IP and http) technology to deliver documents and resources

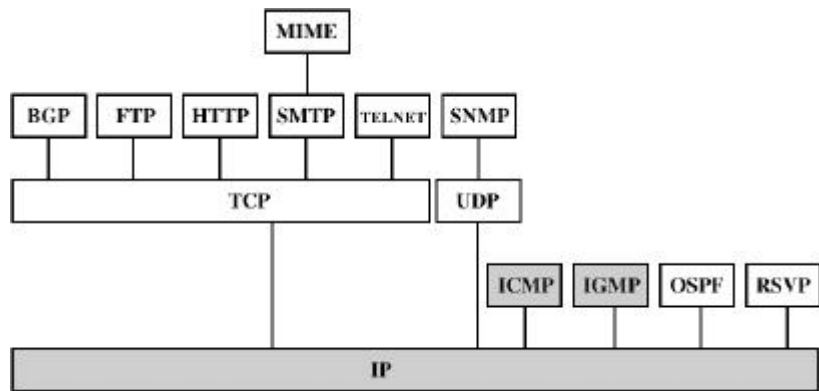
Internetworking Terms (2)

- **End System (ES)**
 - Device attached to one of the networks of an internet
 - Supports end-user applications or services
- **Intermediate System (IS)**
 - Device used to connect two networks
 - Permits communication between end systems attached to different networks

Internetworking Terms (3)

- **Repeater: used to connect link, retransmitting signals bidirectionally - OSI Layer 1**
- **Bridge**
 - IS used to connect two LANs using similar protocols
 - Address filter passing on packets to the required network only
 - OSI layer 2 (Data Link)
- **Router**
 - IS connects two (possibly dissimilar) networks
 - Uses internet protocol present in each router and end system
 - OSI Layer 3 (Network)

Internetworking Protocols



ICMP: internet control message protocol

IGMP: internet group management protocol

Network Architecture Features

- These are generally all present, but different on different constituent networks
 - Addressing
 - Packet size
 - Access mechanism
 - Timeouts
 - Error recovery
 - Status reporting
 - Routing
 - User access control
 - Connection based or connectionless

Architectural Approaches

- Connection oriented
- Connectionless

Connection Oriented

- Assume that each network is connection oriented
- IS connect two or more networks
 - IS appear as DTE to each network
 - Logical connection set up between DTEs
 - Concatenation of logical connections across networks
 - Individual network virtual circuits joined by IS
- May require enhancement of local network services
 - 802, FDDI are datagram services

Connection Oriented IS Functions

- Relaying
- Routing
- e.g. X.75 used to interconnect X.25 packet switched networks
- Connection oriented not often used
 - (IP dominant)

Connectionless Operation

- Corresponds to datagram mechanism in packet switched network
- Each NPDU treated separately
- Network layer protocol common to all DTEs and routers
 - Known generically as the internet protocol
- Internet Protocol (IP)
 - One such internet protocol developed for ARPANET
 - RFC 791
- Lower layer protocol needed to access particular network

Connectionless Internetworking

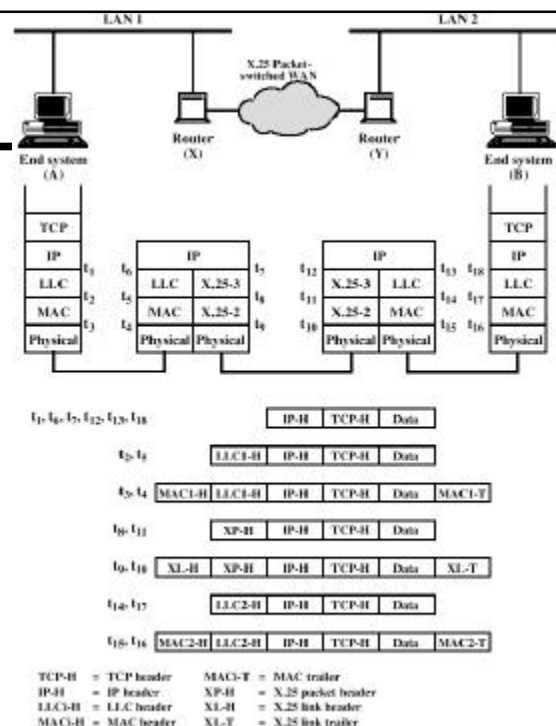
■ Advantages

- Flexibility
- Robust
- No unnecessary overhead

■ Unreliable

- Not guaranteed delivery
- Not guaranteed order of delivery
 - └ Packets can take different routes
- Reliability is responsibility of next layer up (e.g. TCP)

IP Operation



Design Issues

- Routing (already discussed)
- Datagram lifetime
- Fragmentation and re-assembly
- Error control
- Flow control

Datagram Lifetime

- Datagrams could loop indefinitely
 - Consumes resources
 - Transport protocol may need upper bound on datagram life
- Datagram marked with lifetime
 - Time To Live field in IP
 - Once lifetime expires, datagram discarded (not forwarded)
 - Hop count
 - Decrement time to live on passing through a each router
 - Time count
 - Need to know how long since last router
- also important for timeouts at TCP level

Fragmentation and Re-assembly

- Different packet sizes used in different networks
- When to re-assemble
 - At destination
 - | Results in packets getting smaller as data traverses internet
 - Intermediate re-assembly
 - | Need large buffers at routers
 - | Buffers may fill with fragments
 - | All fragments must go through same router
 - Inhibits dynamic routing

IP Fragmentation (1)

- IP re-assembles at destination only
- Uses fields in header
 - Data Unit Identifier (ID)
 - | Identifies end system originated datagram
 - Source and destination address
 - Protocol layer generating data (e.g. TCP)
 - Identification supplied by that layer
 - Data length
 - | Length of user data in octets

IP Fragmentation (2)

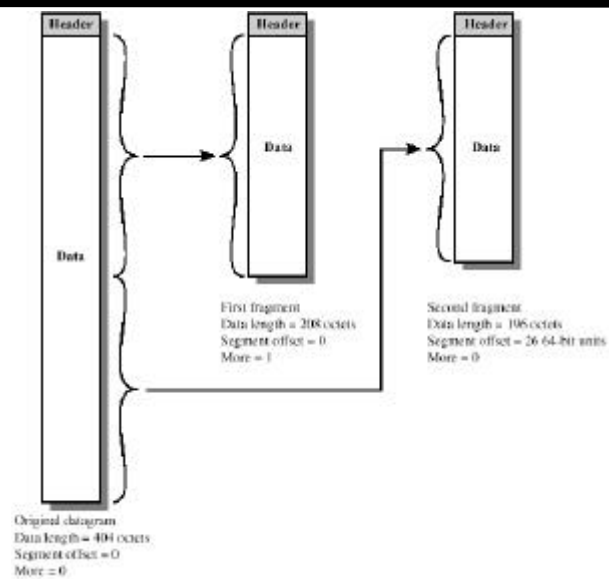
- Offset

- ┆ Position of fragment of user data in original datagram
- ┆ In multiples of 64 bits (8 octets)

- More flag

- ┆ Indicates that this is not the last fragment

Fragmentation Example



Dealing with Failure

- Re-assembly may fail if some fragments get lost
- Need to detect failure
- Re-assembly time out
 - Assigned to first fragment to arrive
 - If timeout expires before all fragments arrive, discard partial data
- Use packet lifetime (time to live in IP)
 - If time to live runs out, kill partial data

Error Control

- Not guaranteed delivery
- Router should attempt to inform source if packet discarded
 - e.g. for time to live expiring
- Source may modify transmission strategy
- May inform high layer protocol
- Datagram identification needed
- (Look up ICMP)

Flow Control

- Allows routers and/or stations to limit rate of incoming data
- Limited in connectionless systems
- Send flow control packets
 - Requesting reduced flow
- e.g. ICMP

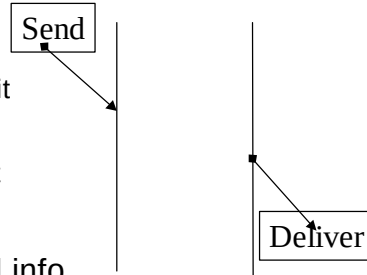
Internet Protocol (IP)

- Part of TCP/IP
 - Used by the Internet
- Specifies interface with higher layer
 - e.g. TCP
- Specifies protocol format and mechanisms

IP Services

■ Primitives

- Functions to be performed
- Form of primitive implementation dependent
 - | e.g. subroutine call
- Send
 - | Request transmission of data unit
- Deliver
 - | Notify user of arrival of data unit



■ Parameters

- Used to pass data and control info

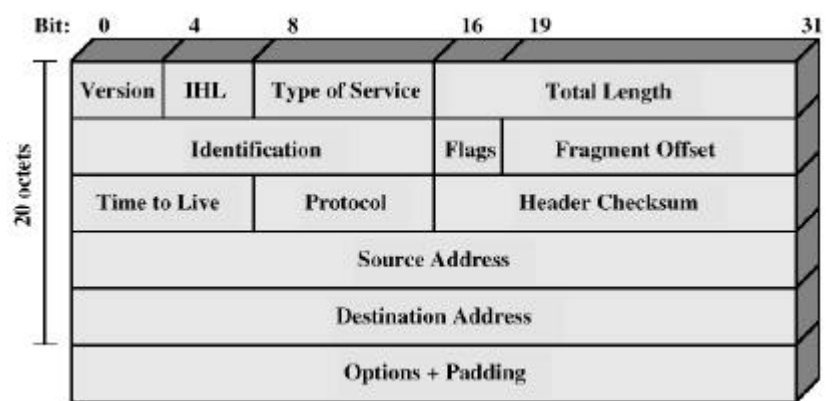
Parameters (1) (Send & Deliver)

- Source address
- Destination address
- Protocol
 - | Recipient e.g. TCP
- Type of Service
 - | Specify treatment of data unit during transmission through networks
 - | Precedence (8 levels), Reliability (Normal or high), Delay (Normal or low), Throughput (Normal or high)
- Identification (Send only)
 - | Uniquely identifies PDU
 - | Needed for re-assembly and error reporting

Parameters (2) (Send & Deliver)

- Don't fragment indicator (Send only)
 - Can IP fragment data ?
 - If not, may not be possible to deliver
- Time to live (Send only)
 - in seconds
- Data length
- Option data
 - security, source routing, route recording, stream id, timestamping
- User data
 - and padding: length is 32 bit multiple

IP Protocol



Header Fields (1)

- Version
 - Currently 4 or 6
- IP header length
 - In 32 bit words
 - Including options
- Type of service
- Total length
 - Of datagram, in octets

Header Fields (2)

- Identification
 - Sequence number
 - Used with addresses and user protocol to identify datagram uniquely
- Flags
 - More bit
 - Don't fragment
- Fragmentation offset
- Time to live
- Protocol
 - Next higher layer to receive data field at destination

Header Fields (3)

- Header checksum
 - Reverified and recomputed at each router
 - 16 bit ones complement sum of all 16 bit words in header
 - Set to zero during calculation
- Source address, Destination address
 - 4 bytes each in IPv4

Data Field

- Carries user data from next layer up
- Integer multiple of 8 bits long (octet)
- Max length of datagram (header plus data)
65,535 octets

IP Addresses - Class A

- 32 bit global internet address
- Network part and host part
- Class A
 - Start with binary 0
 - All 0 reserved
 - 01111111 (127) reserved for loopback
 - Range 1.x.x.x to 126.x.x.x
 - All allocated

IP Addresses - Class B

- Start 10
- Range 128.x.x.x to 191.x.x.x
- Second Octet also included in network address
- $2^{14} = 16,384$ class B addresses
- All allocated

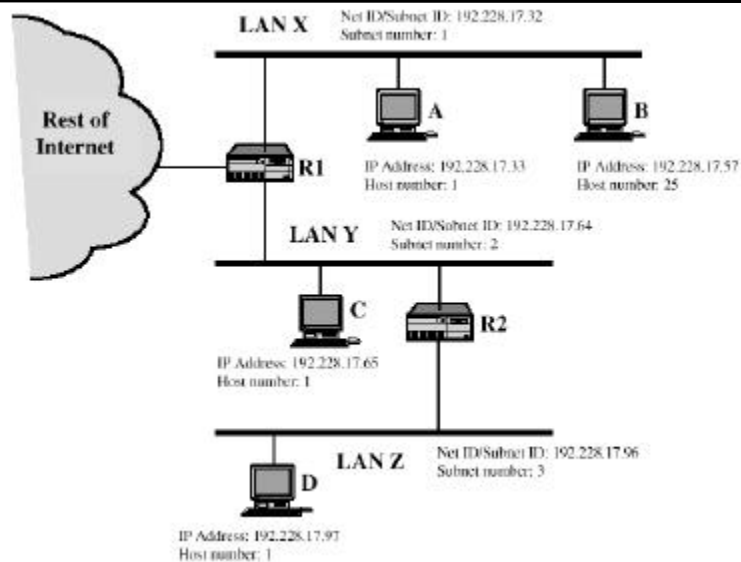
IP Addresses - Class C

- Start 110
- Range 192.x.x.x to 223.x.x.x
- Second and third octet also part of network address
- $2^{21} = 2,097,152$ addresses
- Nearly all allocated
 - See IPv6

Subnets and Subnet Masks

- Allow arbitrary complexity of internetworked LANs within organization
- Insulate overall internet from growth of network numbers and routing complexity
- Site looks to rest of internet like single network
- Each LAN assigned subnet number
- Host portion of address partitioned into subnet number and host number
- Local routers route within subnetted network
- Subnet mask indicates which bits are subnet number and which are host number

Routing Using Subnets



ICMP

- Internet Control Message Protocol
- Transfer of (control) messages from routers and hosts to hosts
- Feedback about problems
 - e.g. destination unreachable, time to live expired, source quench (for flow control)
- Encapsulated in IP datagram
 - Not reliable

ICMP Message Formats

0	8	16	31
Type	Code	Checksum	
Unused			
IP Header + 64 bits of original datagram			

(a) Destination Unreachable; Time Exceeded; Source Quench

0	8	16	31
Type	Code	Checksum	
Pointer	Unused		
IP Header + 64 bits of original datagram			

(b) Parameter Problem

0	8	16	31
Type	Code	Checksum	
Gateway Internet Address			
IP Header + 64 bits of original datagram			

(c) Redirect

0	8	16	31
Type	Code	Checksum	
Identifier		Sequence Number	
Optional data			

(d) Echo, Echo Reply

0	8	16	31
Type	Code	Checksum	
Identifier		Sequence Number	
Originate Timestamp			

(e) Timestamp

0	8	16	31
Type	Code	Checksum	
Identifier		Sequence Number	
Originate Timestamp			
Receive Timestamp			
Transmit Timestamp			

(f) Timestamp Reply

0	8	16	31
Type	Code	Checksum	
Identifier	Sequence Number		

(g) Address Mask Request

0	8	16	31
Type	Code	Checksum	
Identifier		Sequence Number	
Address Mask			

(h) Address Mask Reply